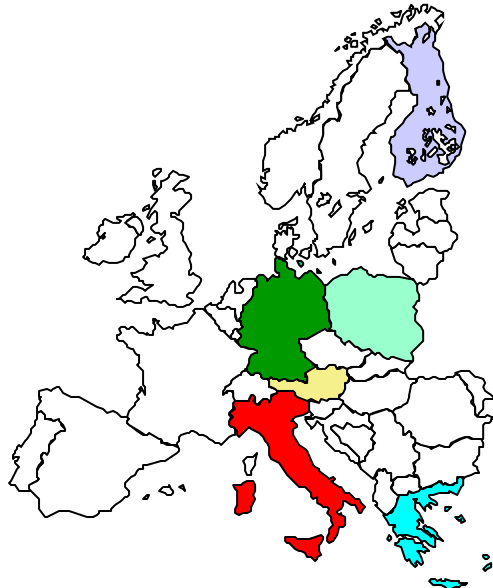
**AQUILA (IST-1999-10077)**



**Adaptive Resource Control for QoS
Using an IP-based Layered Architecture**



**A QoS Architecture with
Adaptive Resource Control:
The AQUILA Approach**



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<http://www.ist-aquila.org/>

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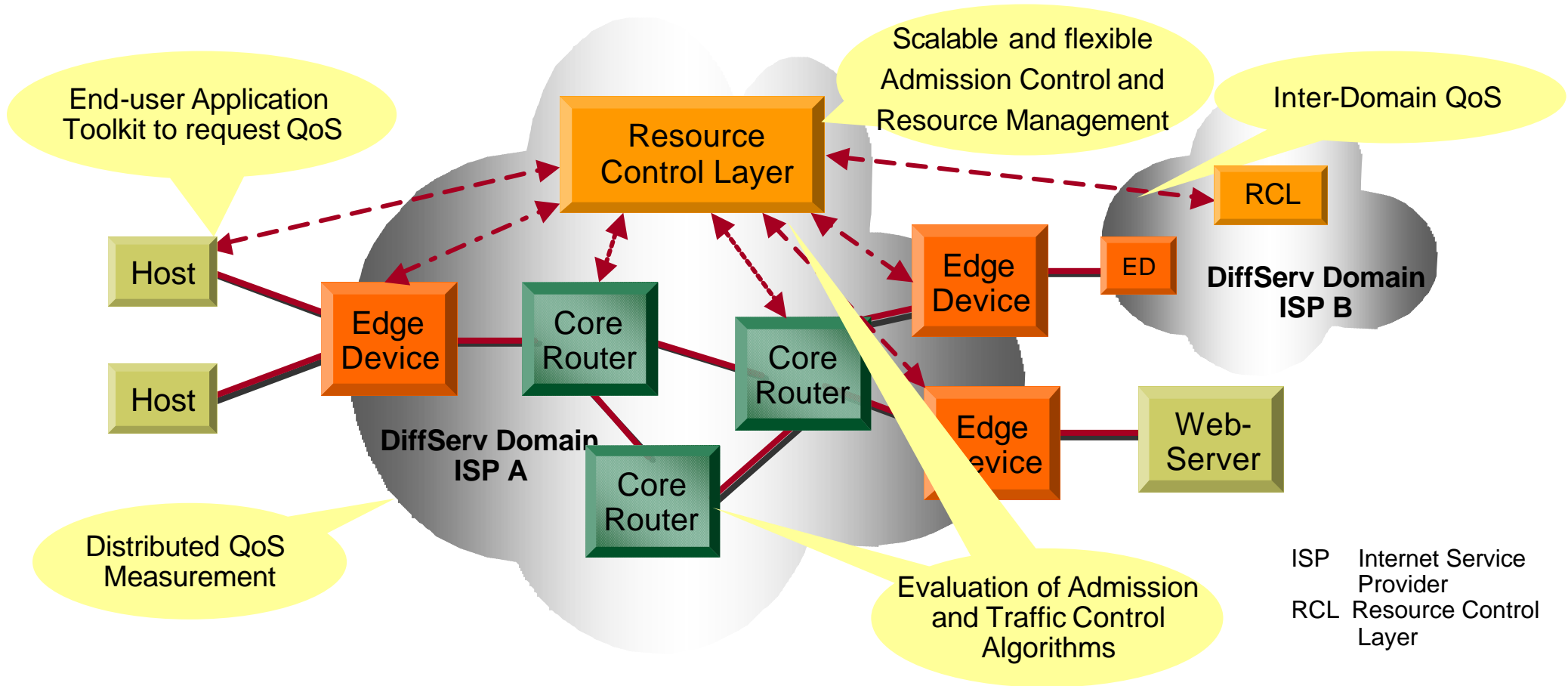


Warsaw Univ. of Technology,
Poland

Main Objectives

- Investigate dynamic end-to-end QoS Provisioning in IP Networks
- Implement Prototypes of a QoS Architecture for a Carrier Grade DiffServ Core Network
- Support a wide Range of Applications by providing a QoS Toolkit / API
- Continuously analyse Customer Requirements, Market Situations and Technological Trends and develop Business Models
- Contribute to Standardisation Bodies like IETF, ITU, ETSI, etc.

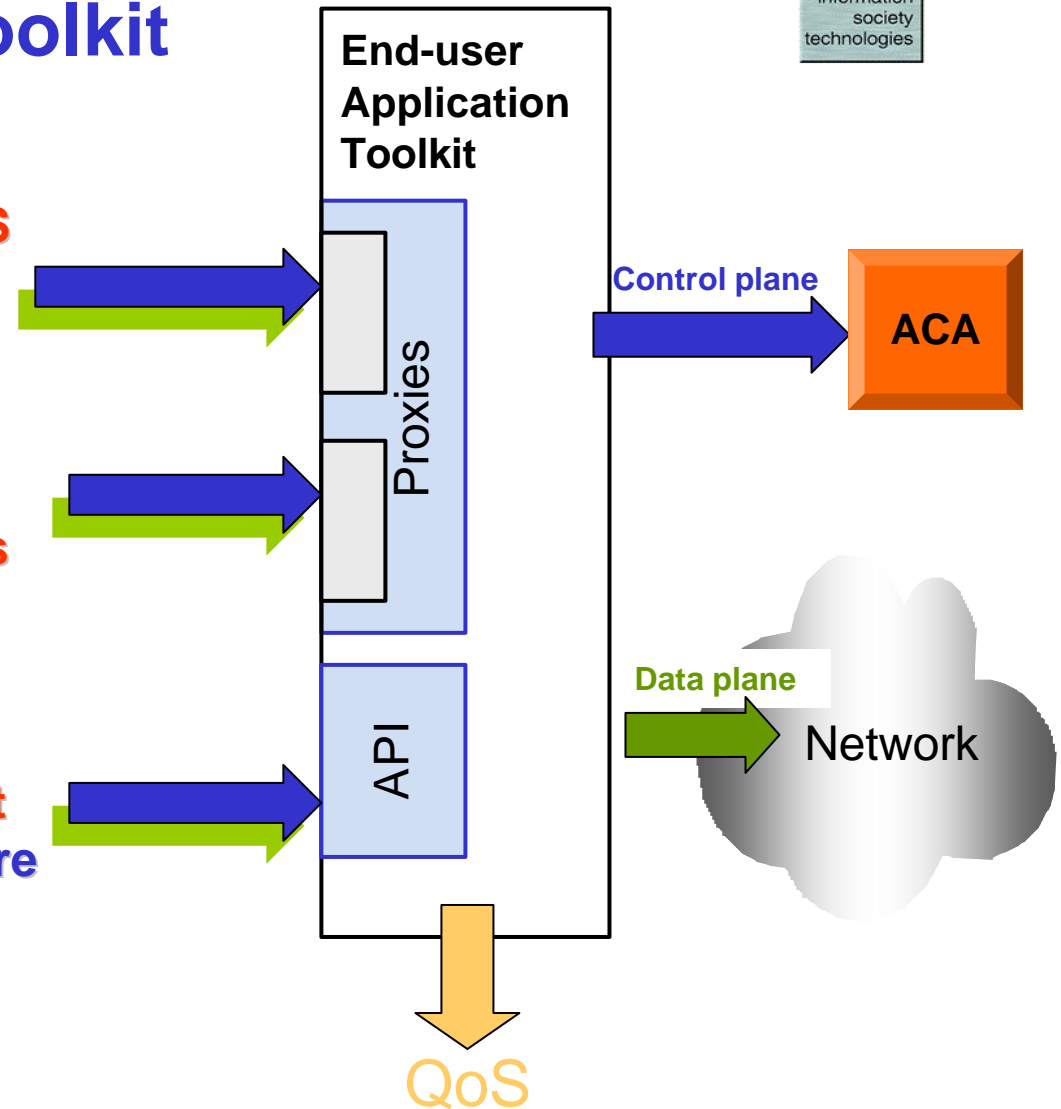
Main Innovations



end-to-end Quality of Service

End-user Application Toolkit Objectives

- Enable Access to QoS to non QoS aware Legacy Applications
- Support QoS aware Applications (RSVP, DiffServ) / Support various QoS Request Methods
- Provide a Methodology and a Programming Interface to support the Construction of new QoS aware Applications
- Provide an End-user friendly QoS Access



Design Goals

■ Scalable Architecture

- Distributed building blocks
- Autonomous operation of elements

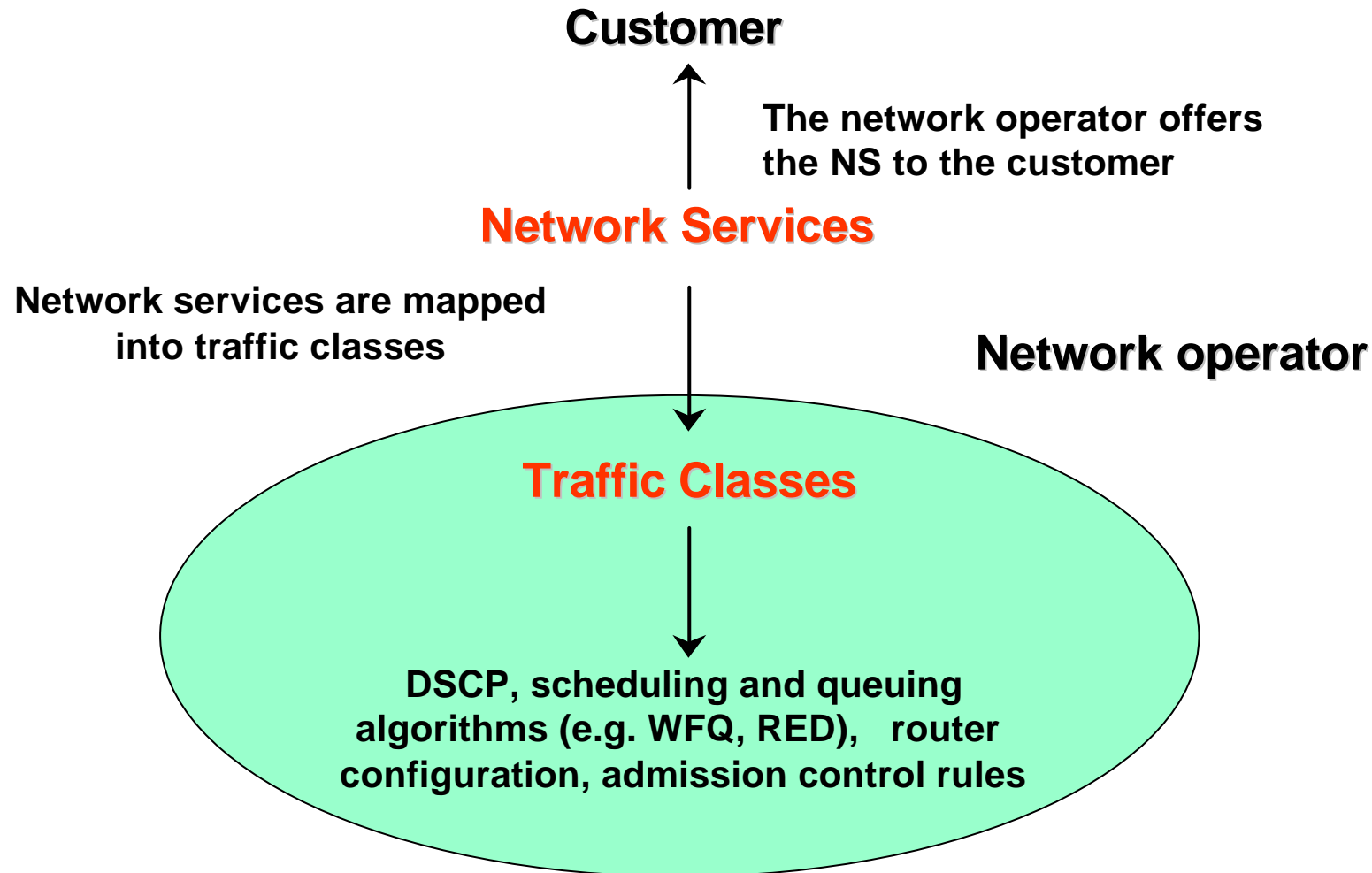
■ Failure Proof

- Failure of an element should only degrade (if at all), not disable the operation of other elements

■ Based on DiffServ Core Network

- Use of existing, commercial routers
- Enable migration path from current networks

Traffic Classes

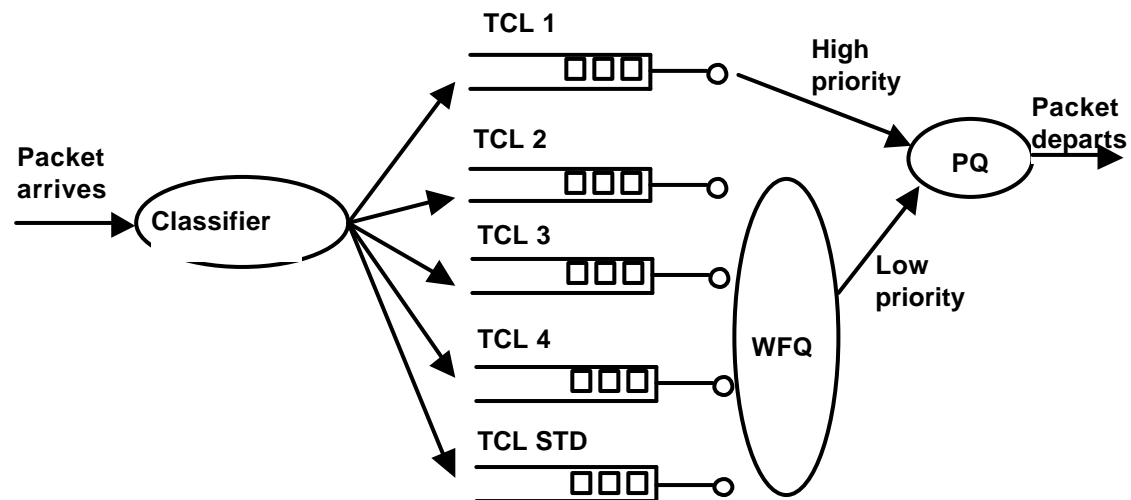


Traffic Classes

■ Five Traffic Classes have been specified

Network service	Premium CBR	Premium VBR	Premium MultiMedia	Premium Mission Critical	Standard
Traffic class	TCL 1	TCL 2	TCL 3	TCL 4	TCL STD

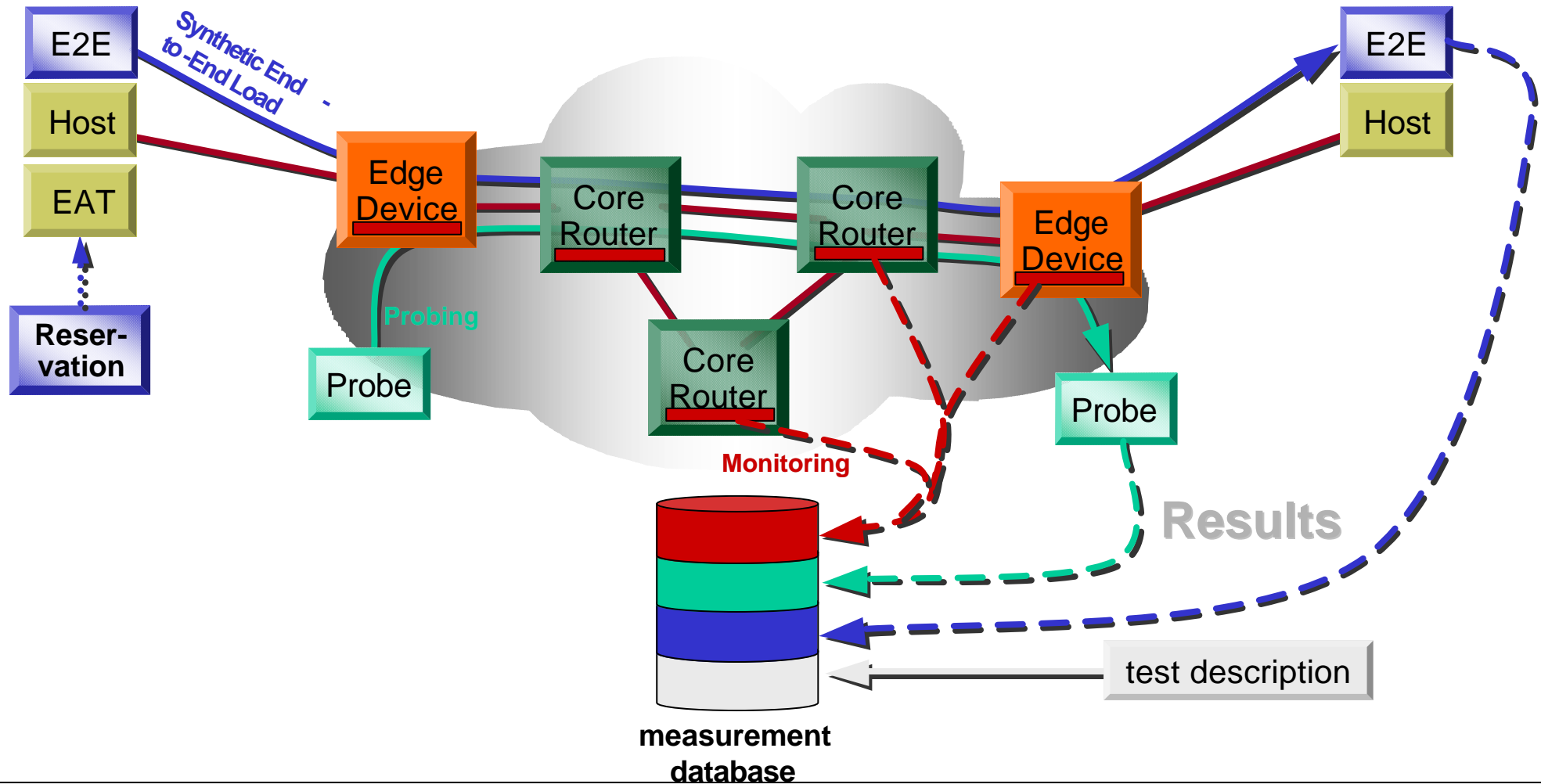
■ ... as well as the related Traffic Control Mechanisms in the Routers



Distributed QoS Measurements Packet Level

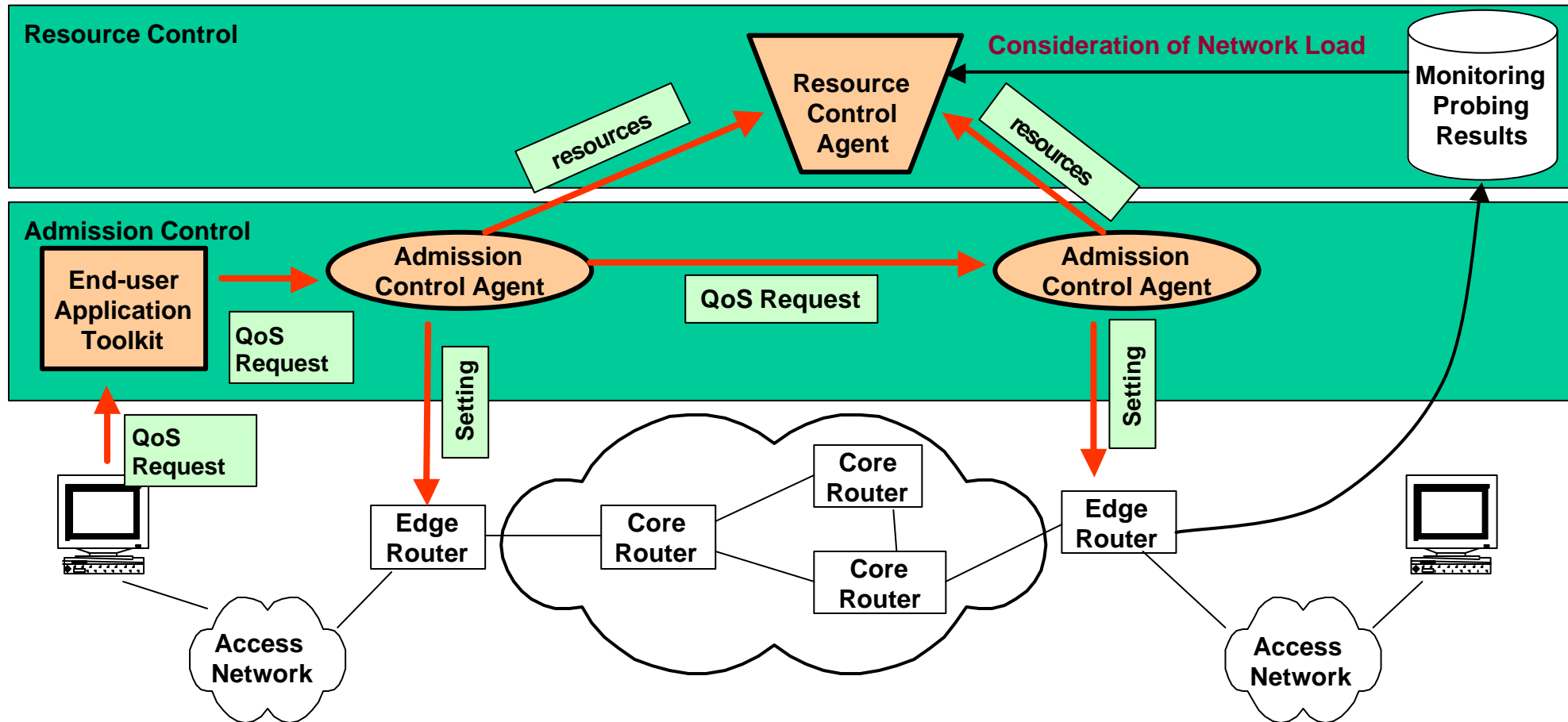
- **Measurement of end-to-end Parameters via Probes**
 - Examples: one-way delay, delay variation, packet loss
- **Collection of Performance Monitoring and QoS Parameters from Routers**
- **Definition and Implementation of Synthetic Flow Load Generators for Measurements of end-to-end QoS**
- **Storage of all Measurement Data in Measurement Database**

Measurements



Architecture

Resource Control Layer



Resource Control Layer (1)

■ Basic Idea of DiffServ Network

- Provide some (fixed) prioritised traffic classes within the network
- Guarantee QoS by limiting amount of prioritised traffic at the network edge (limited resources)

■ Additional Benefit of the Resource Control Layer

- Dynamically shift resources between network edges (P resource pools)
- Take into account the actual resource usage of the network (2nd Trial)

Resource Control Layer (2)

■ Admission Control Agent

- Authenticates user
- Authorises and checks request
- Locates ingress and/or egress edge router (P roles)
- Requests resources from the resource control agent
- Admits / rejects new flows
- Installs policies in ingress router

■ Resource Control Agent

- Manages resources
- Checks availability of requested resources
- (Re-)distributes resources as needed

Resource Pools

■ Resource Limits

- Limit amount of QoS traffic from each edge router

■ Group neighboured Routers

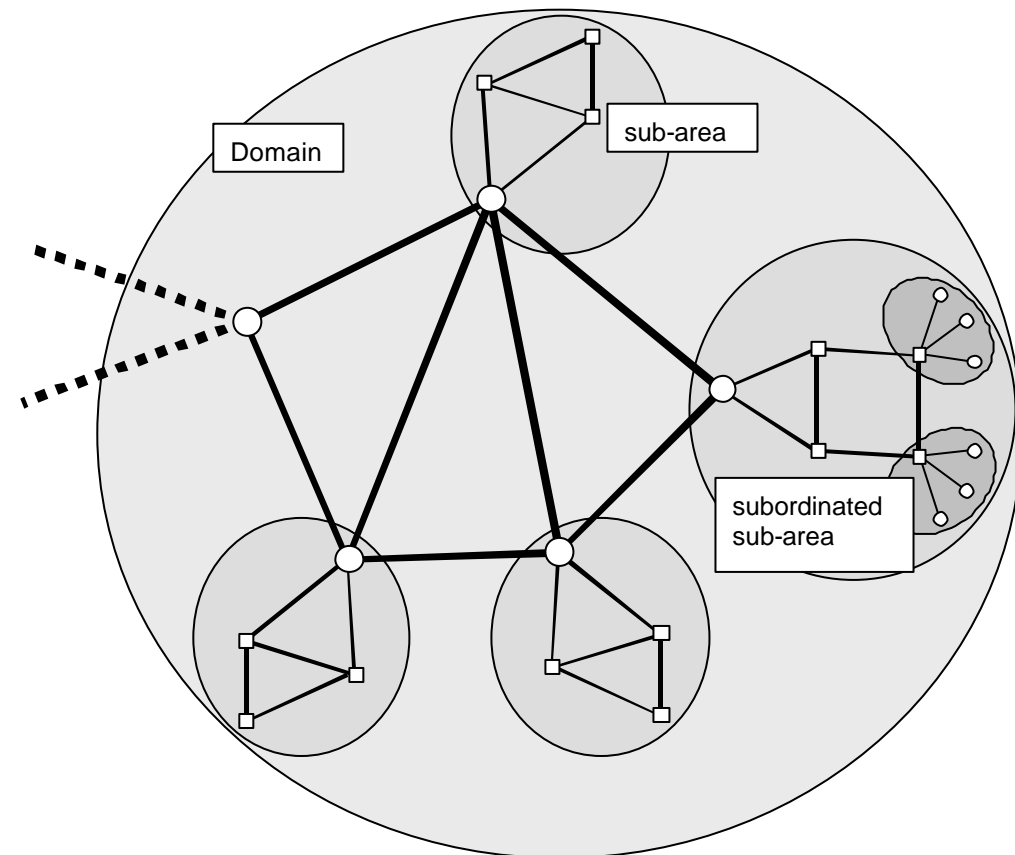
- Limit amount of QoS traffic from each group

■ Dynamic Distribution

- Dynamically shift resources within group

■ Hierarchical Structure

- “Groups of groups”



Results

■ Network services

- fulfils assumptions
 - Delay, packet loss \leq specification
- do not mix streaming and elastic traffic in one service
 - else UDP will degrade TCP && UDP will not reach QoS
- Fair bandwidth sharing for TCP flows

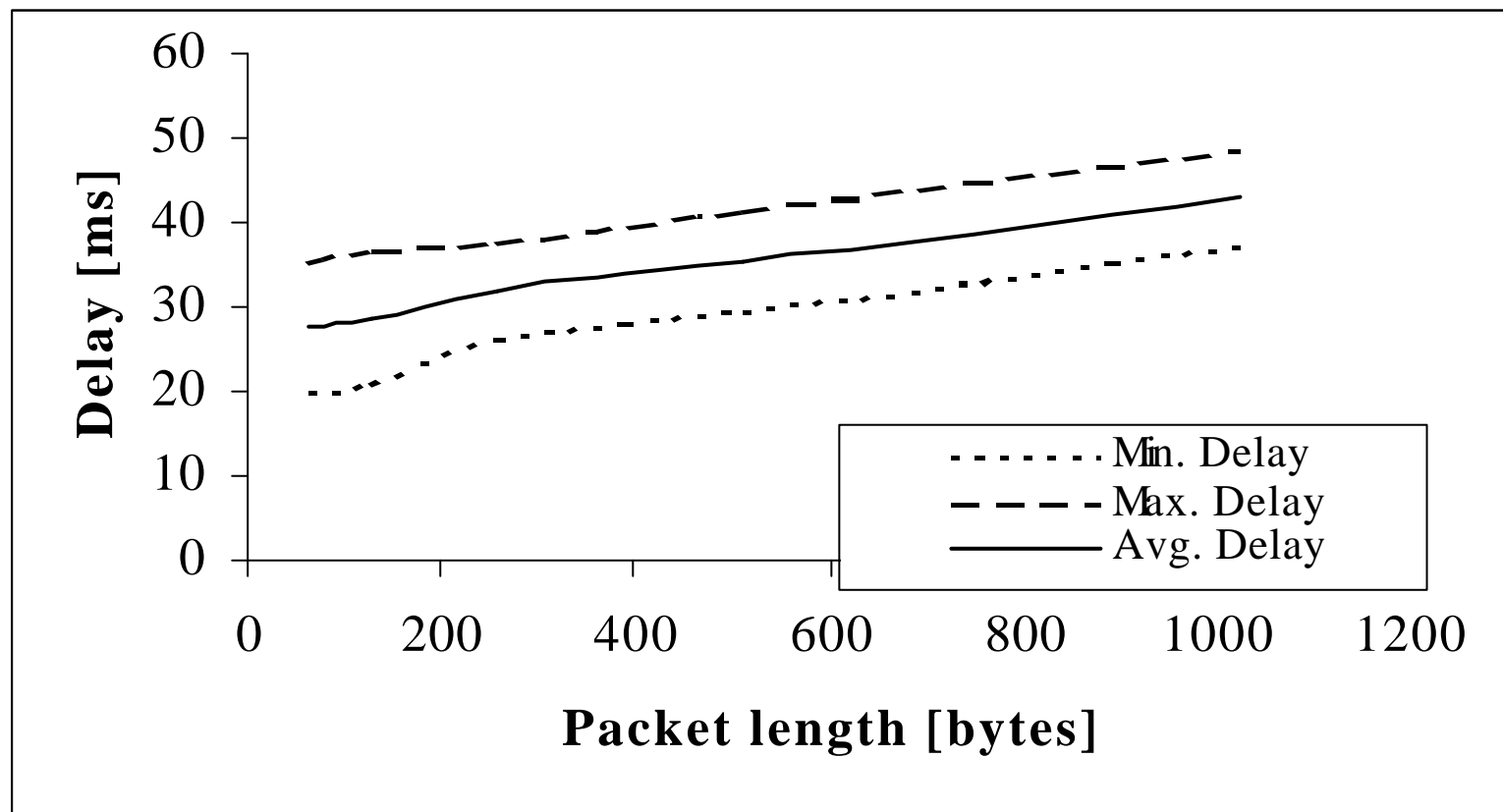
■ Resource and admission control

- General mechanisms work
 - still more evaluations and simulations (large networks) necessary
- Ongoing work on effective bandwidth calculations
 - not always correct and not manageable at the moment

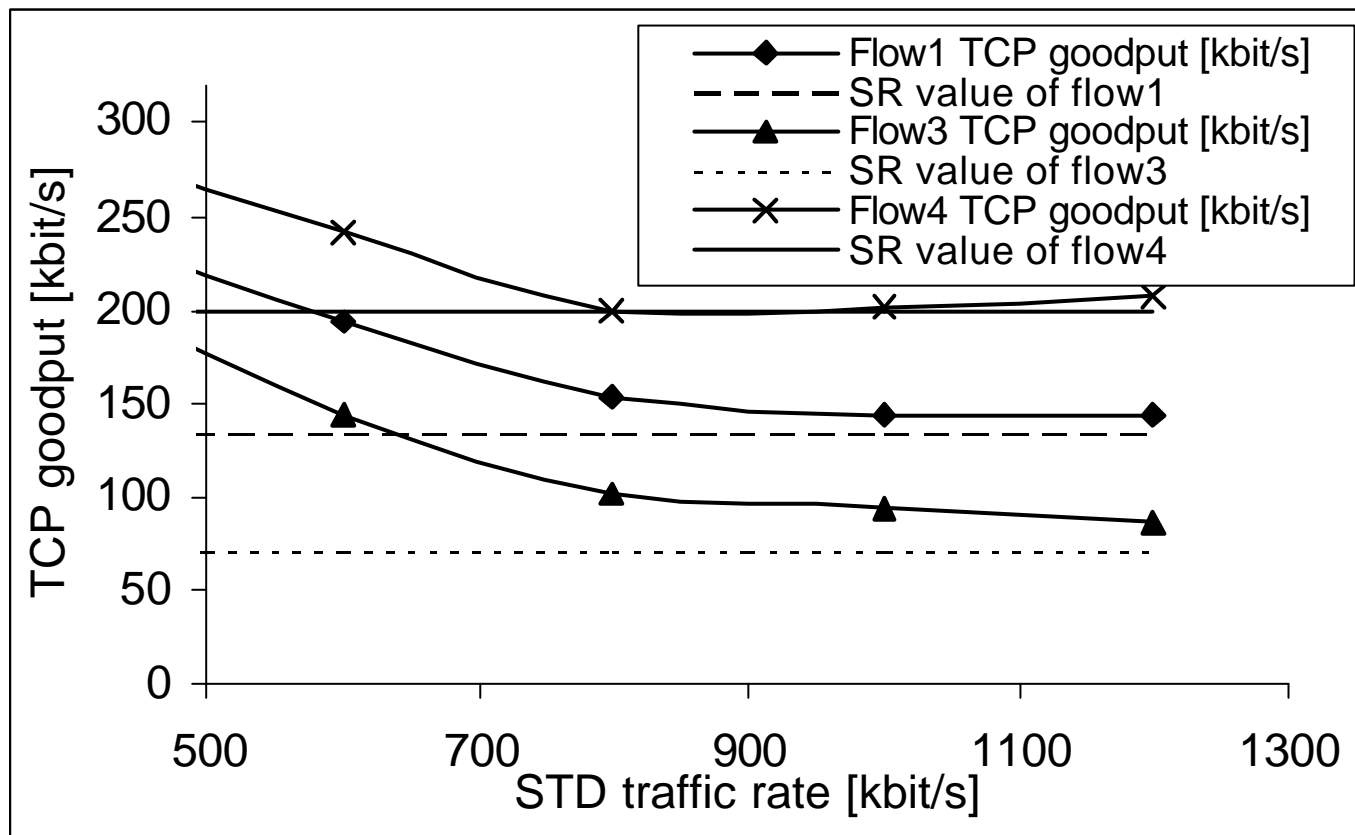
■ Router

- WFQ noticeably degrades performance of Cisco routers

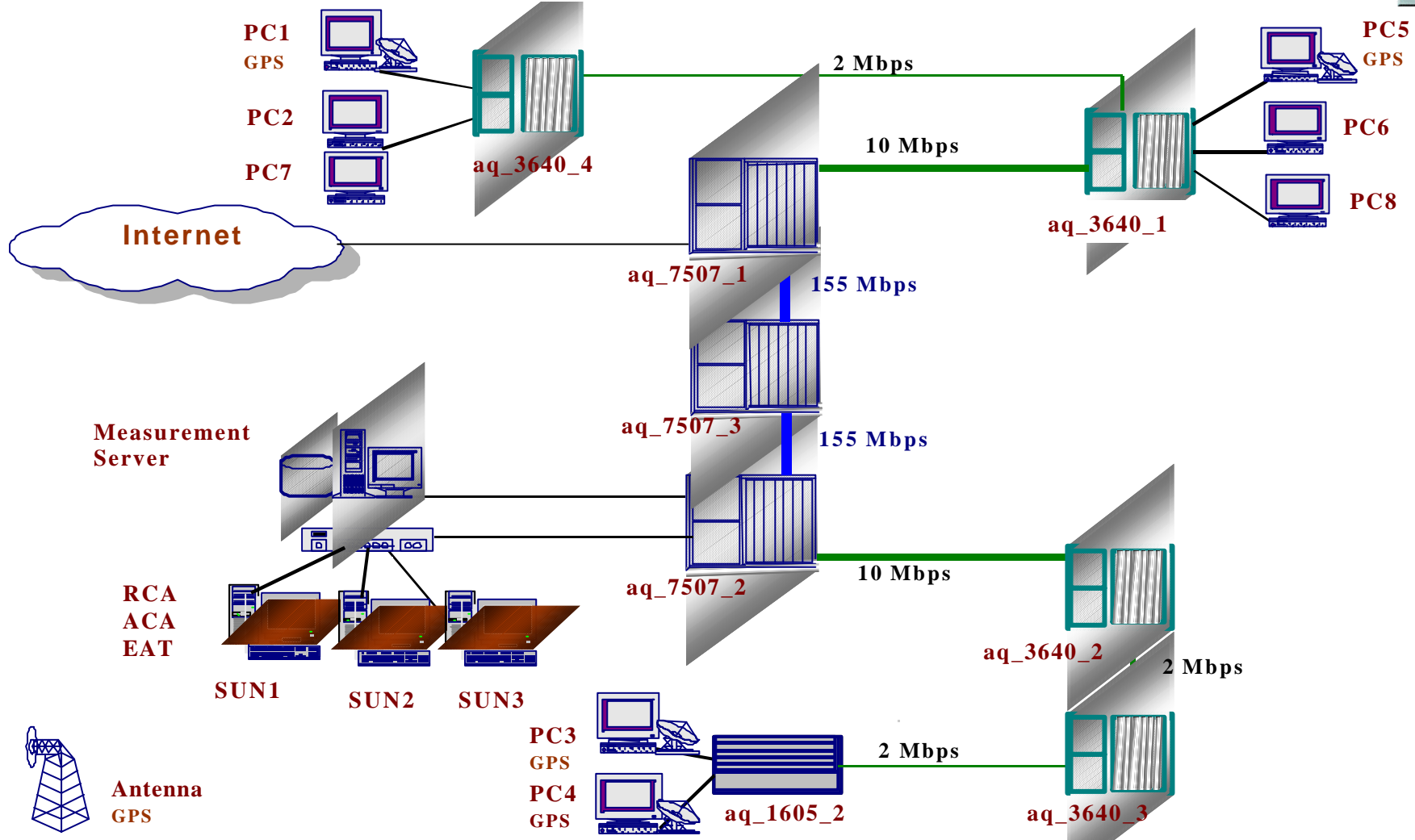
PCBR: Delay experiment



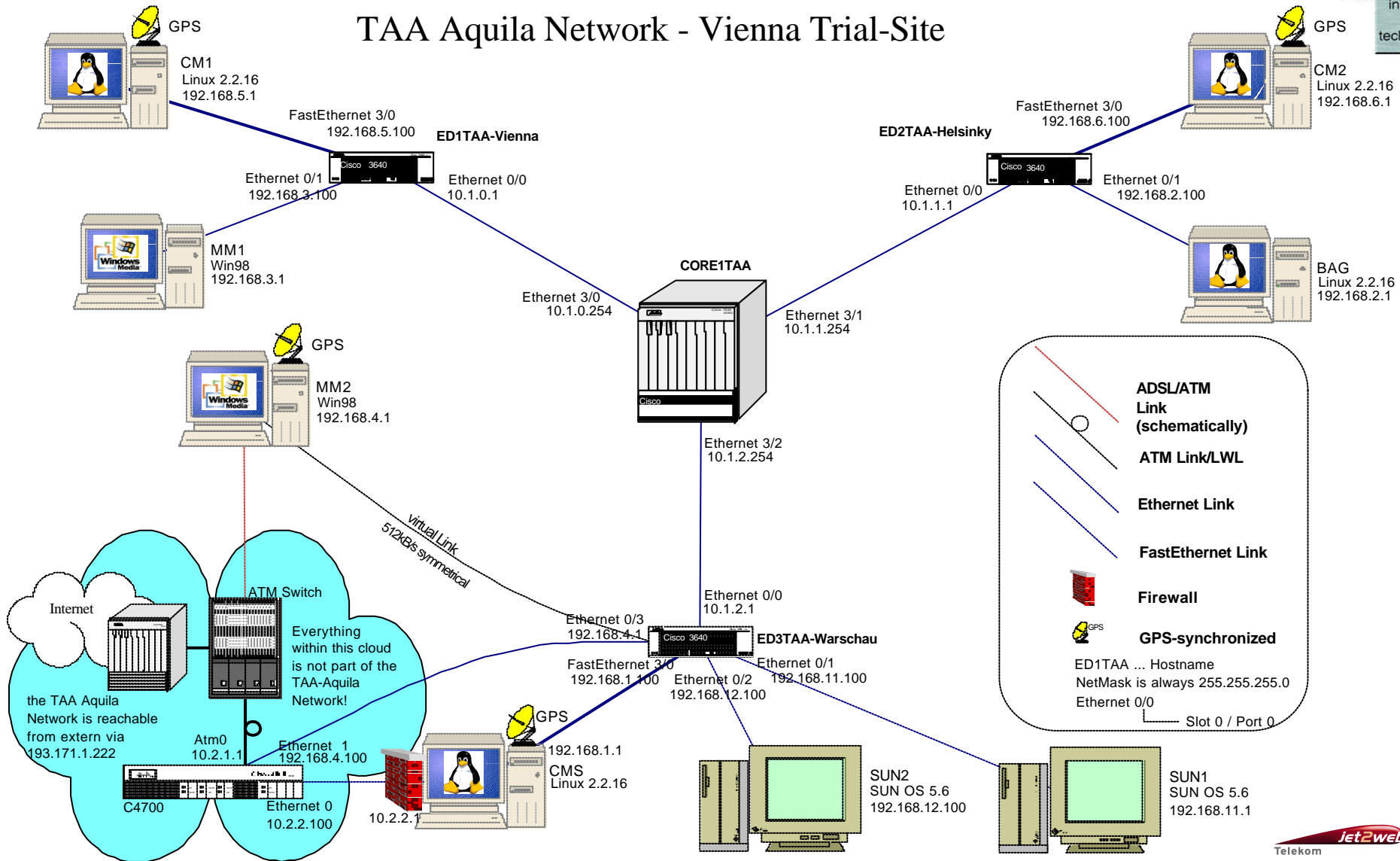
Goodput of TCP flows with different SR



Warsaw Testbed

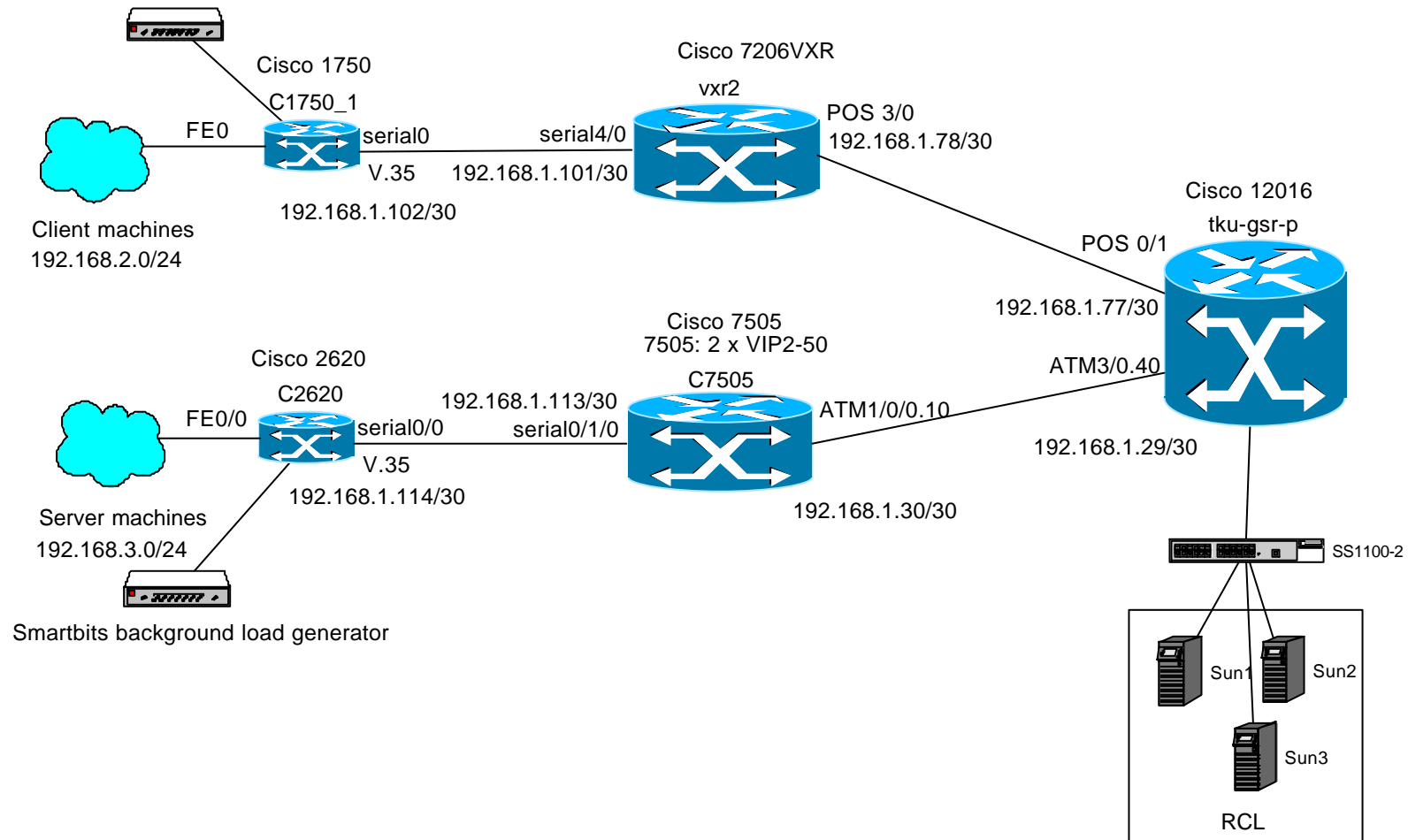


TAA Aquila Network - Vienna Trial-Site



Helsinki Testbed (Elisa Communications)

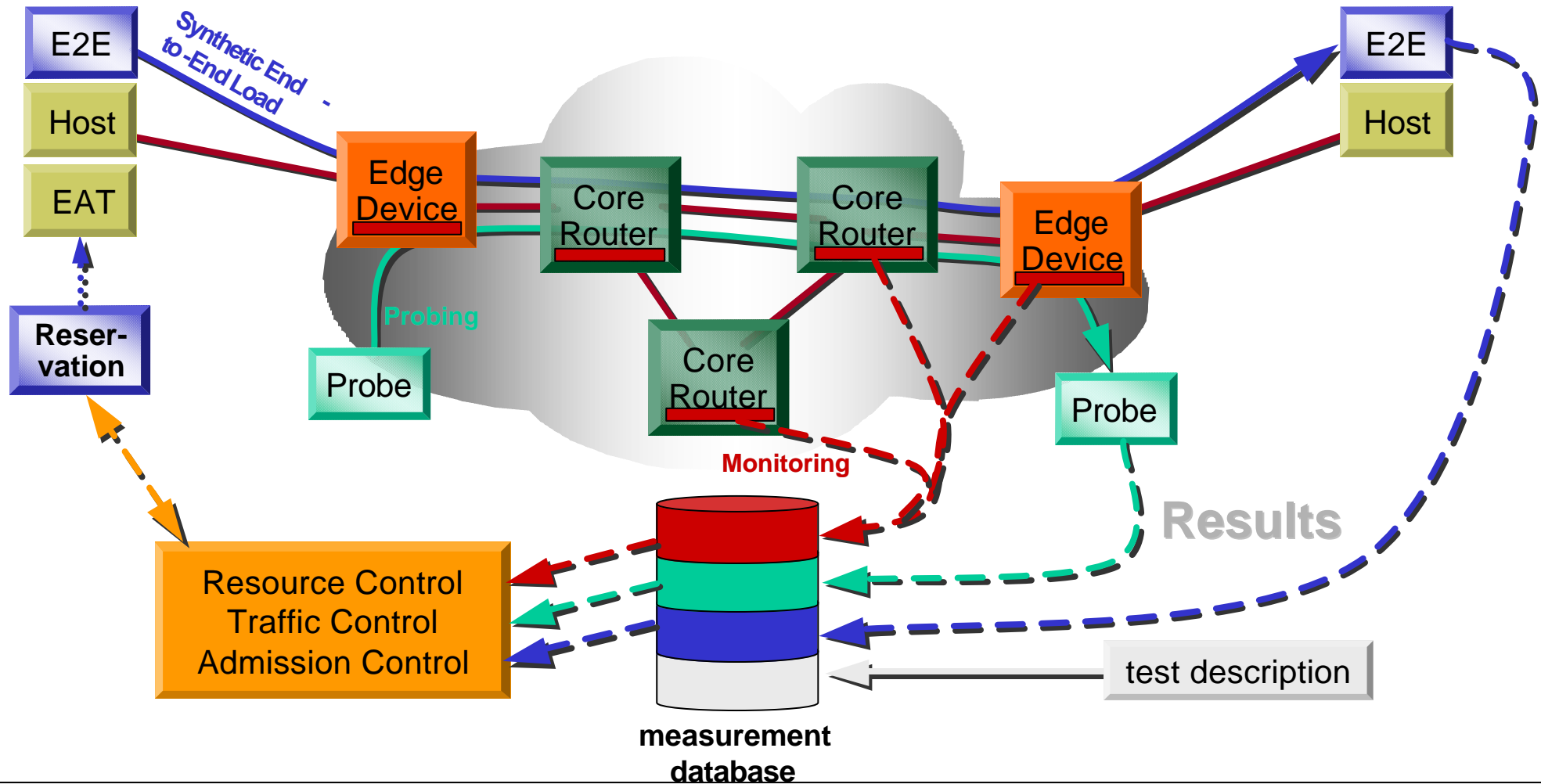
Smartbits background load generator



Future work

- **Control loops**
 - for provisioning
 - for admission control
- **Inter-domain resource allocation**
- **Management tool**
- **Interoperation with MPLS**
- **SIP protocol gateway**

Control loops (1)



Control loops (2)

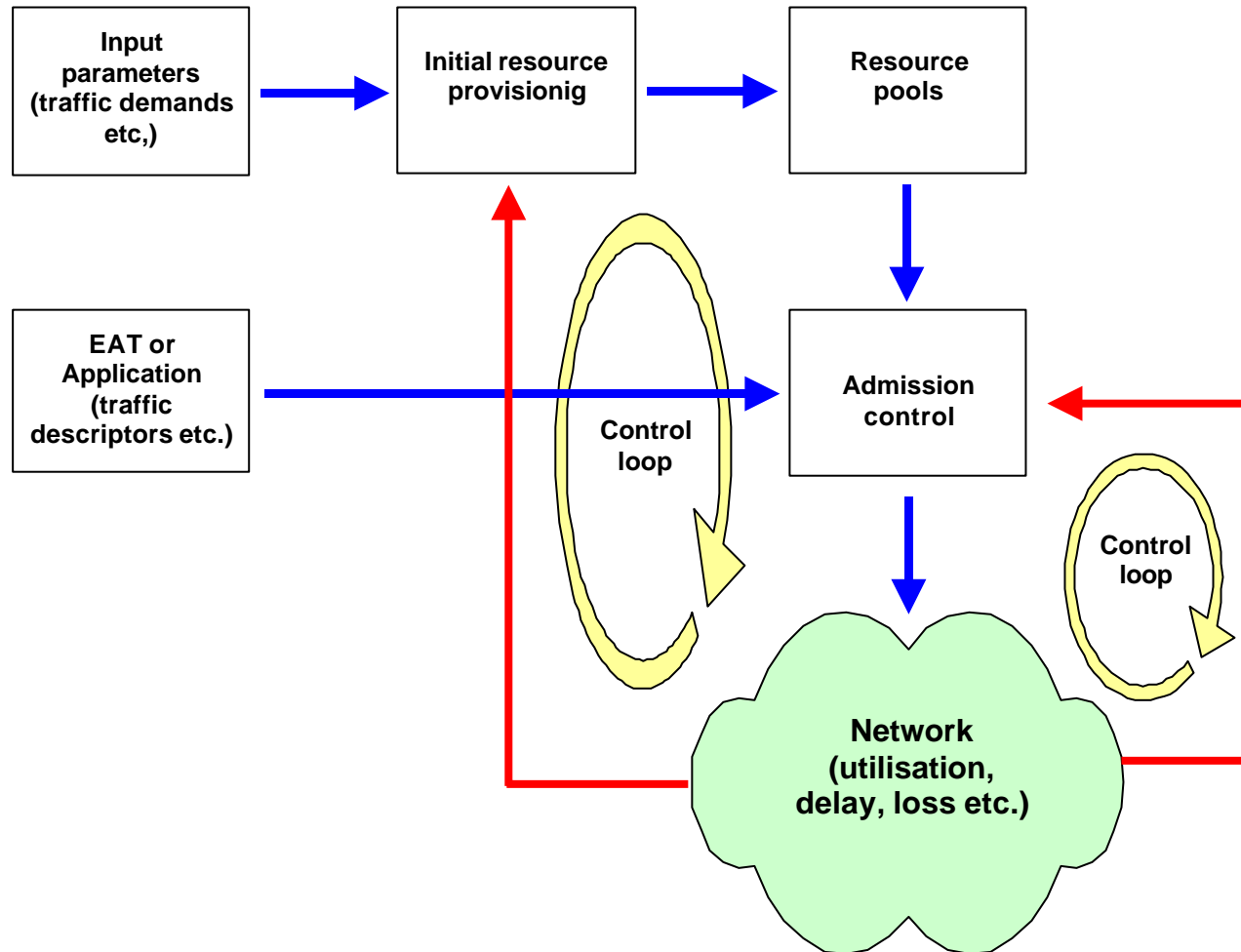
■ Open loop

- Resource distribution on the basis of a priori specified parameters
- User declarations for admission control
- Assumed traffic demand for resource provisioning

■ Closed loop

- Resource distribution on the basis of parameters obtained from the network
- Feedback from measurements
- Resource control loop
- Admission control loop (MBAC)

Control loops (3)



Internet domain architecture

■ Domains, autonomous systems (AS)

- Currently, the Internet consists of about 10.000 AS
- Each AS is operated and managed by a network operator

■ Internet routing

- Internally, each AS may use any routing protocol
- For routing between AS, BGP (Border Gateway Protocol) is used world-wide

■ Independent, but co-operating routing mechanisms

- Same is required for resource allocation
- Inter-domain resource control architecture must not depend on any particular resource control mechanism within the AS

BGRP: Tackling the scalability problem

■ Microflow reservation

- At a core router, there might be more than 1 Mio individual flows at the same time

■ AS (autonomous system) pair reservation

- We could aggregate reservations starting and ending at the same AS pair into one inter-domain reservation
- This would still yield in the order of 10.000 reservations on a core router at the same time
- In general, the problem remains, that the number of reservations grows $O(N^2)$ with the number of AS in the internet

BGRP: Sink tree based aggregation

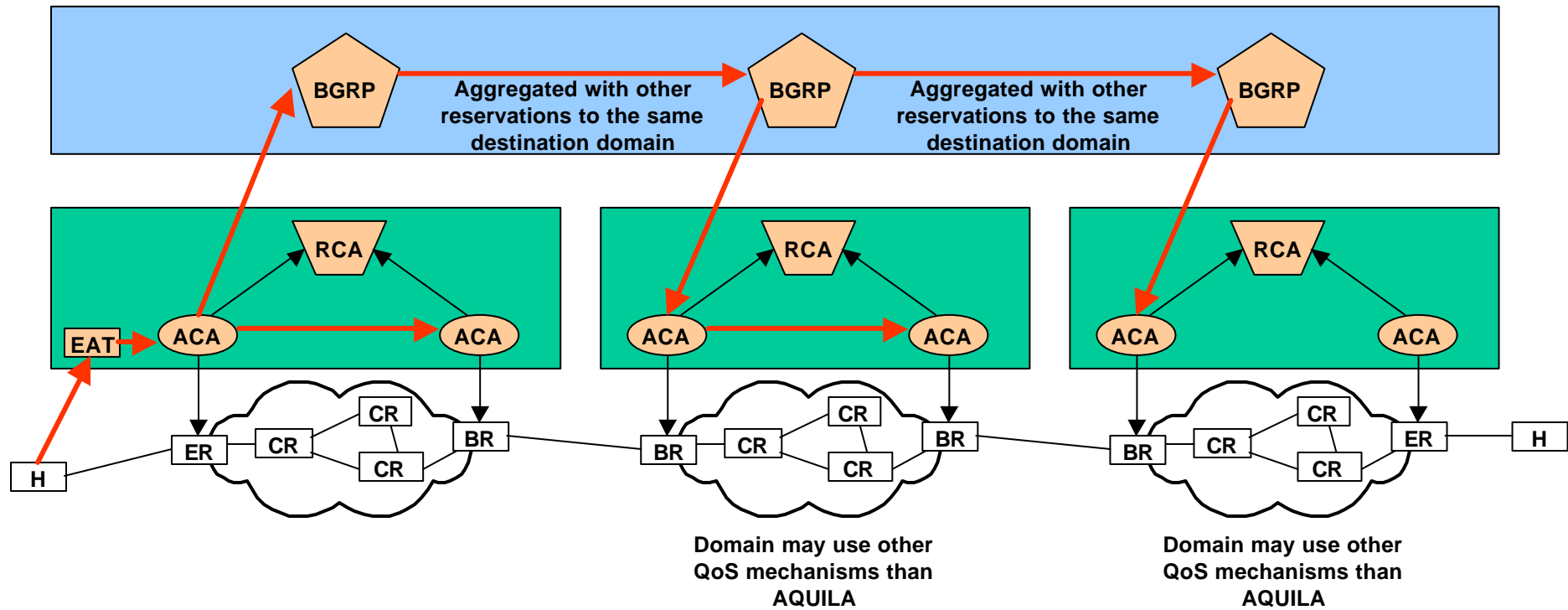
■ Sink tree: definition of term

- An BGP router sends all traffic for the same destination AS to the same next hop AS (property of the BGP routing protocol)
- This guarantees the construction of a so called sink tree for a destination AS
- The root of the sink tree is the destination AS
- The traffic from all other AS travels along the links of this sink tree to the destination AS

■ Aggregation

- Reservations from various AS to a common destination AS can be aggregated, as they merge along the sink tree

Inter-domain Resource Control



BGRP: Relation to intra-domain reservations

■ Operates on the boundaries between ISPs

- Each ISP is free to operate its own domain independently
- Additional layer above the intra-domain resource control

■ Relation to AQUILA phase 1

- AQUILA phase 1 ACA requests inter-domain resources from the BGRP bandwidth broker at the initiating domain
- BGRP bandwidth broker requests intra-domain resources from the AQUILA phase 1 architecture at subsequent domains

Comparison: Sink trees vs. resource pools

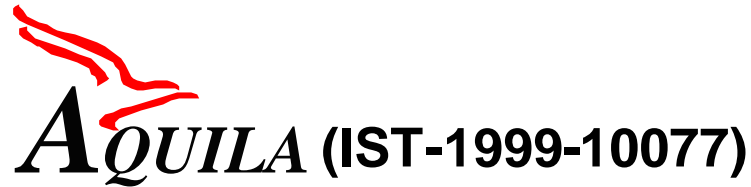
■ Intra-domain advantages of resource pools

- Resource pools can be used for ingress and egress admission control
- Resource pools scale for networks with many sources and sinks (Scales better than $O(\text{hosts})$)
- Configuration and administration based on an overall picture of the domain.

■ Inter-domain advantages of BGRP sink trees

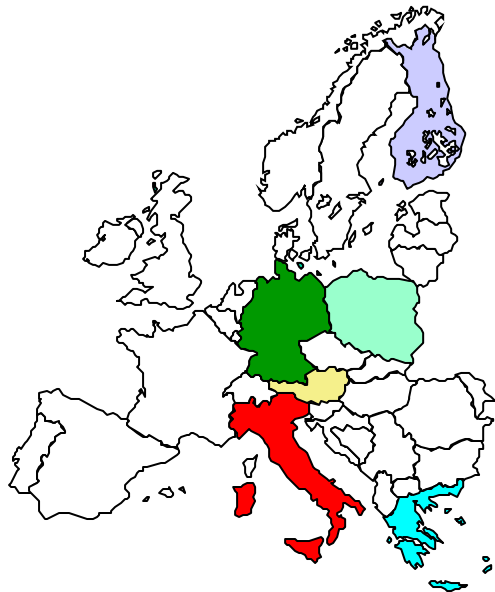
- BGRP considers the full network topology at the AS level
- BGRP can be independently configured and administered at each AS

→ Both approaches are well positioned within an overall AQUILA architecture

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**Thank you for
your attention !**

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